

Computability and Logic

dBerLog

Q1 2007

Bachelor Information Meeting

- Department of Computer Science offers a meeting for students briefing on the bachelor year and the master's studies
- Time: 15-17, November 1, 2007
- Place: Store Auditorium

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- Summary
 - Universality
 - Duality
 - Self-reference
 - Program Verification
- Life stories
- About the exam

Universal Computational Models

Turing: Turing machines (1930s)

Goedel: recursive functions (1930s)

Church: λ -calculus (1930s)

Chomsky: Language grammars (1950s)

All these have been shown to be “equivalent” wrt expressiveness!

Church's λ -calculus

- $(\lambda x. E) (F) \xrightarrow{\beta} E [F/x]$
- $(\lambda x. E) \xrightarrow{\alpha} \lambda y. (E [y/x])$

- **Theorem**

A (partial) function is computable in the λ -calculus iff it is computable by a Turing machine!

Chomsky type 0 grammars

- Context dependent rules:
 $\alpha A \beta \rightarrow \alpha \gamma \beta$ where $A \in V$ and $\alpha, \beta, \gamma \in (V \cup T)^*$

- **Theorem**

The class of languages generated by Chomsky type 0 grammars is exactly the class of (Turing) recursively enumerable languages!

Gödel's μ -recursive functions $N^k \rightarrow N$

- Successor, zero-test, projections, function composition, and primitive recursion:

$$f(0, x) = h(x)$$

$$f(n+1, x) = g(n, x, f(n, x))$$

unbounded minimization:

$$f(n) = \min y. (g(y, n) = 0)$$

- **Theorem**

A (partial) function is definable as a μ -recursive function iff it is computable by a Turing machine.

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Hækleopskrift: en dug...

Opskrift/program:

Hækl 21 lm

1. række: * 1 stm, 1 lm, spring over næste lm *; Gentag fra * til * 9 gange til; 1 stm i sidste lm (10 mlmrum), vend
- 2.-10. række: * 1 stm på stm, 1 lm, spring over mlmrum *; Gentag fra * til * 9 gange til;

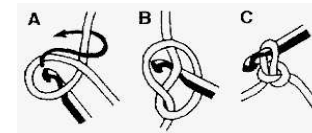
.....

Materialer/data:

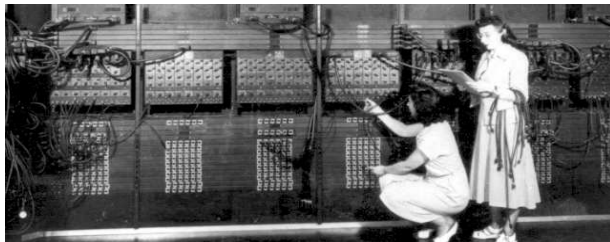
Perlebomuld # 5
Hæklenål 1,9 mm

...og en hækler

En luftmaske



Programming ENIAC 1943-45



Duality between programs and data

John von Neumann
1903-1957

First draft on a report on
EDVAC 1945



Duality between programs and data

Alan Turing
1912-1954

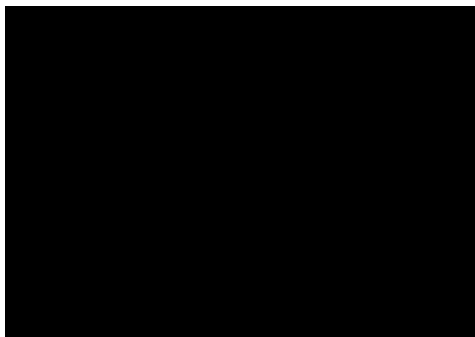
On Computable Numbers
with an application to the
Entscheidungsproblem
1936



Bootstrapping



Duality Rene Magritte (1966)

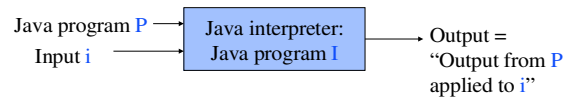


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Programs applied to programs

Happens ALL the time! Example: a Java interpreter



Text/program manipulation - examples

- Encryption programs
- Program optimizers
- Indentation programs

Proof reading

SIGNS USED IN CORRECTING PROOFS		SIGNS USED IN CORRECTING PROOFS	
Δ Delete; take out	\times Change broken letter	Δ Delete; take out	\times Change broken letter
\odot Close up	Stet Let it stand as set	\odot Close up	Stet Let it stand as set
\wedge Insert	$\dots\dots$ Let it stand as set	\wedge Insert	$\dots\dots$ Let it stand as set
$\#$ Insert space	wf Wrong font, size or style	$\#$ Insert space	wf Wrong font, size or style
\lrcorner Raise	$l.c.$ Lower case, not capitals	\lrcorner Raise	$l.c.$ Lower case, not capitals
\llcorner Lower	$rom.$ Use Roman letter	\llcorner Lower	$rom.$ Use Roman letter
\square Move to left	blf Use black type letters	\square Move to left	blf Use black type letters
\square Move to right	\odot Period	\square Move to right	\odot Period
\parallel Straighten type line at side of page	Δ Comma	\parallel Straighten type line at side of page	Δ Comma
\parallel Straighten lines	\sphericalangle Apostrophe	\parallel Straighten lines	\sphericalangle Apostrophe
$\text{\textcircled{P}}$ Paragraph	$\text{\textcircled{S}}$ Superior figure	$\text{\textcircled{P}}$ Paragraph	$\text{\textcircled{S}}$ Superior figure
center Put in middle of page or line	$\text{\textcircled{I}}$ Inferior figure	center Put in middle of page or line	$\text{\textcircled{I}}$ Inferior figure
$\text{\textcircled{H}}$ Transpose	$\text{\textcircled{H}}$ Hyphen	$\text{\textcircled{H}}$ Transpose	$\text{\textcircled{H}}$ Hyphen
$\text{\textcircled{h}}$ Transpose	$\text{\textcircled{h}}$ Use small capitals	$\text{\textcircled{h}}$ Transpose	$\text{\textcircled{h}}$ Use small capitals
$\text{\textcircled{c}}$ Turn inverted letter right side up	$\text{\textcircled{c}}$ Use capitals	$\text{\textcircled{c}}$ Turn inverted letter right side up	$\text{\textcircled{c}}$ Use capitals
	$\text{\textcircled{i}}$ Use italics		$\text{\textcircled{i}}$ Use italics

Self-reference

All books in our library have a list of references

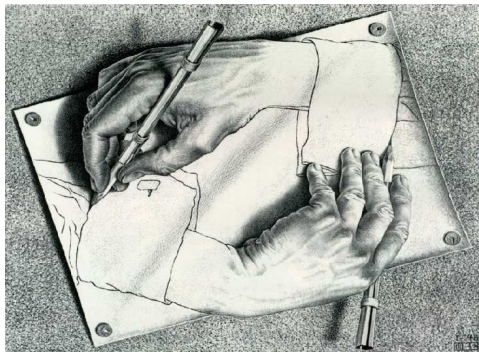
Some books reference themselves (“see chapter..”)

Let us call such a book self-referencing

Task:

write a book (for the library) containing a list of all the books (in the library), which are not self-referencing!

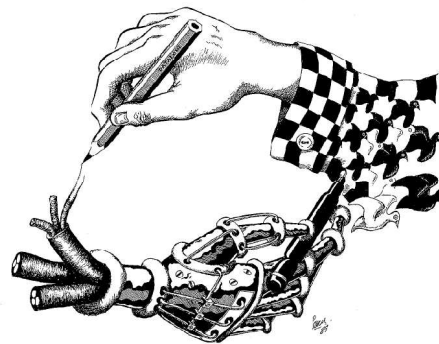
Selfreference - Escher 1898-1972



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Selfreference - dBerLog



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Program correctness

Algorithm: Euklid (m, n)
Inputbetingelse: $m, n \geq 1$
Outputkrav: $r = \text{sfd}(m, n)$
Metode: $\{m, n \geq 1\}$
 $p \leftarrow m; q \leftarrow n;$
 $I = \{\text{sfd}(p, q) = \text{sfd}(m, n)\}$ while $p \neq q$ do
 if $p > q$ then $p \leftarrow p - q$
 else $q \leftarrow q - p;$
 $r \leftarrow p$
 $\{r = \text{sfd}(m, n)\}$



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Program incorrectness

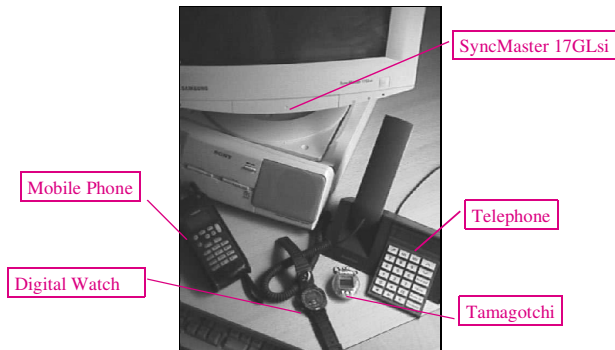
Algorithm: Euklid (m, n)
Inputbetingelse: $m, n \geq 1$
Outputkrav: $r = \text{sfd}(m, n)$
Metode: $\{m, n \geq 1\}$
 $p \leftarrow m; q \leftarrow n;$
 $I = \{\text{sfd}(p, q) = \text{sfd}(m, n)\}$ while $p \neq q$ do
 if $p > q$ then $p \leftarrow p - q$
 else $q \leftarrow q - p;$
 $r \leftarrow q$
 $\{r = \text{sfd}(m, n)\}$



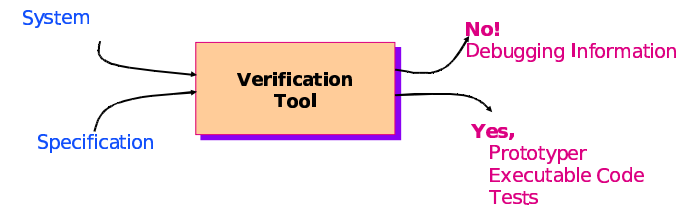
Program termination

```
{x > 1}
while x ≠ 1 do
  if even(x) then x := x div 2 else x := 3 × x + 1
{true}
```

Embedded Systems



Automated Verification



A REAL system



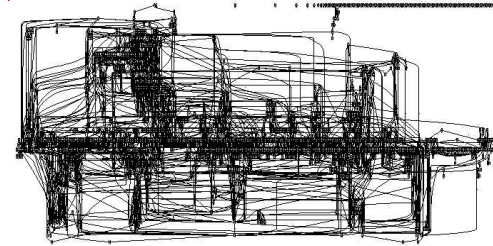
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Train Simulator

1421 machines
11102 transitions
2981 inputs
2667 outputs
3204 local states
Declare state sp.: 10^{476}

BUGS ?



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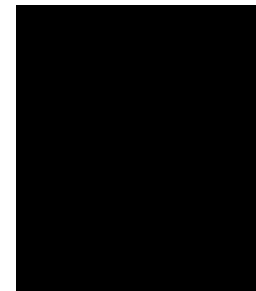
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Life stories - A. Turing 1912-1956



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Life stories - K. Gödel 1906-1978

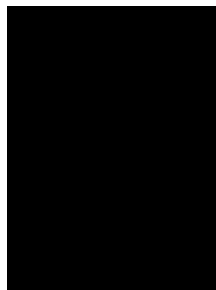


Ud og se med DSB - september 2007

Ax. 1. $\bullet \forall x \{ [P(x) \rightarrow Q(x)] \wedge P(y) \rightarrow P(y) \}$
Ax. 2. $P(x) \leftrightarrow \neg \neg P(x)$
Th. 1. $P(x) \rightarrow \exists x P(x)$
Ef. 1. $G(x) \rightarrow \forall y [P(y) \rightarrow P(x)]$
Ax. 3. $P(G)$
Th. 2. $\exists x G(x)$
Th. 2. $\varphi \leftrightarrow \psi \leftrightarrow \varphi(x) \wedge \forall y [\varphi(y) \rightarrow \psi(y)]$
Ax. 4. $P(x) \rightarrow \bullet P(x)$
Th. 3. $G(x) \rightarrow G(x)$
Ef. 3. $E(x) \rightarrow \forall y [P(x) \rightarrow \bullet P(y)]$
Ax. 5. $P(E)$
Th. 4. $\bullet \exists x G(x)$

KURT GÖDELS GUDSBEVIS
Gödel indviede kun sine nærmeste venner i sit gudsbevis, for han mente, at det havde visse svagheder. Først efter hans død blev det publiceret. Gödels gudsbevis er en matematisk version af teodiske teologer som Augustin af Caesarees gudsbevis, som går på følgende: Det hører med til gudsarbejdet, at Gud er det mest fuldkomne væsen. Man kan derfor ikke forestille sig noget, der er mere fuldkomment end Gud. Derfor må Gud eksistere, for hvis han ikke eksisterede, ville han mangde eksistens, og mangende eksistens ville være en ufuldkommenhed ved Gud. Hvis Gud altså ikke eksisterede, kunne man forestille sig et endnu mere fuldkomment væsen, som så over at ligne Gud også eksisterede. Men det ville sige mod gudsarbejdet, for man kan nemlig forestille sig noget mere fuldkomment end Gud. Ergo må Gud eksistere.

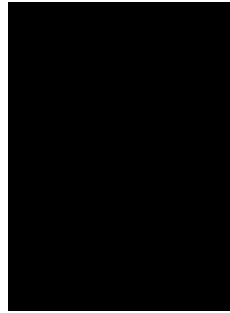
Life stories - E. Post 1897 - 1958



Life stories - A. Church 1903 - 1995



Life stories - N. Chomsky 1928 -



Life stories - T. Hoare 1934 -



Life stories - J. C. Martin ?? -



Life stories - J. Kelly ?? - 1995

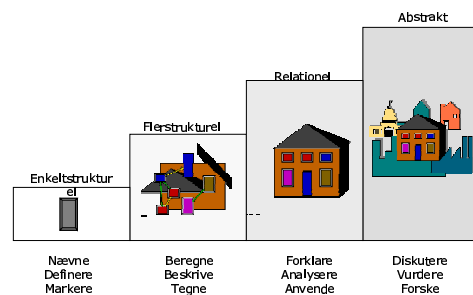
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 - Diagonalization
 - Self-reference
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dBerLog - Exam

- **Oral exam**
 - Friday October 19 - Wednesday October 24 2007
 - Internal examiners:
 - Anders Møller, Michael Schwartzbach, Ole Østerby
 - Two questions:
 - Computability, Logic
 - 20 minutes each - no preparation time
- **Compulsory home works**
 - 2 compulsory written home work assignments must all be handed in and accepted by the tutor
 - OBS STRICT DEADLINE TOMORROW!!!!!!

SOLO taxonomi



dBerLog - Goals

- The goals of this course are to give the student the following capabilities
 - to be *familiar* with the basic *terminology* for computability and logic
 - to *describe* basic computability classes and fundamental logics
 - to *describe* basic *properties* of computability classes and logics
 - to *explain* constructive/algorithmic approaches to computability classes and logics
 - to *analyse* and to *prove* properties of computability classes and logics

dBerLog - Goals Computability

- The goals of this course are to give the student the following capabilities
 - to *be familiar with* the basic *terminology* for computability
 - problems as formal languages and operations on these, decidability, Turing machines
 - to *describe* basic computability classes and their properties
 - recursive and recursively enumerable languages, closure and decidability properties, from intuition and examples to formal notation and definitions
 - to *explain* algorithmic approaches to properties of computability classes
 - constructive arguments for closure and decidability properties, problem reductions
 - to *analyse* and to *prove* properties properties of computability classes
 - diagonalization, reduction

dBerLog - Goals Logic

- The goals of this course are to give the student the following capabilities
 - to *be familiar with* the basic *terminology* for logic
 - truth, satisfaction, validity, syntax, semantics
 - to *describe* fundamental logics and their properties
 - propositional logic, truth tables, predicate logic, interpretations and valuations, program logics, proof systems
 - to *explain* algorithmic approaches to properties of logics
 - decidability, normal forms, proofsystems and their proofs, from examples to formal definitions
 - to *analyse* and to *prove* properties properties of logics
 - soundness and completeness, existence and non-existence of proof systems, Gödel's theorems

Plans for the 7 weeks

- Model of Computation: Turing Machines
- Computability
- Non computable problems

- Propositional Logic
- Predicate Logic
- Program Logic - Gödel's theorems

- Summary - Exam

dBerLog Curriculum

Martin:

chapter 9, chapters 10.1-10.2, 10.3, 10.5 (excl. proofs of Thms 10.8 and 10.9), chapter 11 (excl. proof of Thms 11.14 and 11.15)

Kelly:

chapter 1, chapter 4, chapter 6.1-6.7.4, 6.9-6.10, chapter 7

Nielsen:

Limitations of Program Verification, 2004

Turing machines

- *Definition and operations of TMs* - examples of TMs solving problems, computing functions
- *Variations of TMs*
- *Churh-Turing thesis*
- *The universal TM*

Recursive and recursively enumerable languages

- *Definitions of RE and R*
- *Closure properties* of RE and R
- *Characterizations of RE* - enumerating a language, Chomsky grammars,...
- *Countability arguments* for the existence of non-RE languages

Unsolvable problems

- *Definition of solvable problem*
- *The languages NSA and SA* - and their membership wrt. RE and R
- *The reduction technique*
- *Unsolvable problems for TMs*
- *Rice's theorem*
- *Other unsolvable problems*, - PCP and CFG problems

Logic - semantics

- Clear understanding of logical *syntax* and *semantics*, logical *truth*, *satisfaction*, *validity*, logical *connectives*, logical *consequence*, *models*
- Propositional logic: *truth tables*
- Predicate logic: *variables*, *objects*, *predicates*, *functions*, *quantifiers*, *scope*, *binding*, *substitution*, - *interpretations*, *valuations*

Logic proof systems

- General definition of *axiomatic proof system* and its theorems
- Proof systems *AL* for propositional logic and *FOPL* for predicate logic in particular
- *Soundness* and *completeness* of axiomatic proof systems
- *Proofs* of soundness and completeness of *AL*, *decidability* for propositional logic
- *Awareness* of soundness, completeness, decidability results for *FOPL*/predicate logic

Limitations of program verification

- *Hoare triples, partial and total correctness*
- *Hoare proof system*, - and its relation to proof system for the model of natural numbers
- *Incompleteness theorem for Hoare triples* - and its proof: provability rec. enum. - truth not rec. enum.
- *Goedels incompleteness theorem* - and its proof (see above)

dBerLog exam

- Please make sure all formalities are in place - registration, compulsory exercises, etc!
- Please show up for your exam well in advance of your scheduled time !
- And remember to enjoy the exam.....